7. Behavioural Modelling

David Pereira José Proença Eduardo Tovar

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Requirements and Model-driven Engineering

CISTER – ISEP
Porto, Portugal

https://cister-labs.github.io/ramde2122

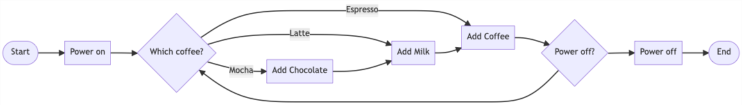
Overview

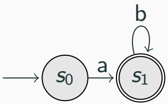
So far

- Models and properties for structures: boolean and 1st order logic, ...
- Useful, e.g., for UML class diagrams

Next

- Look at UML behaviour diagrams
- Use a domain with a precise semantics
  - Non-deterministic finite automata (NFA)
  - Simple language for processes
  - Encode processes $\rightarrow$ NFA
  - Equivalence of processes





What are formal methods?

Formal methods are techniques to

model complex systems using

rigorous mathematical models

Specification

Define part of the system using a modelling language

Verification

Prove properties.

Show correctness.

Find bugs.

Implementation

Generate correct code.

All formal models are wrong

All formal models are wrong

... but some of them are usefull!

Syllabus

- High-level overview or requirements and associated processes
- Mathematical Preliminaries
  - Basic mathematical notations
  - Set theory
  - PropositionalLogic
  - First Order Logic
- Behavioural modelling
  - Single component
    - State diagrams and Flow charts
    - Formal modelling: Automata, Process Algebra in mCRL2
  - Many components
    - Communication diagrams and Sequence diagrams
    - Formal modelling: Process algebra with interactions
  - Equivalences
  - Verification

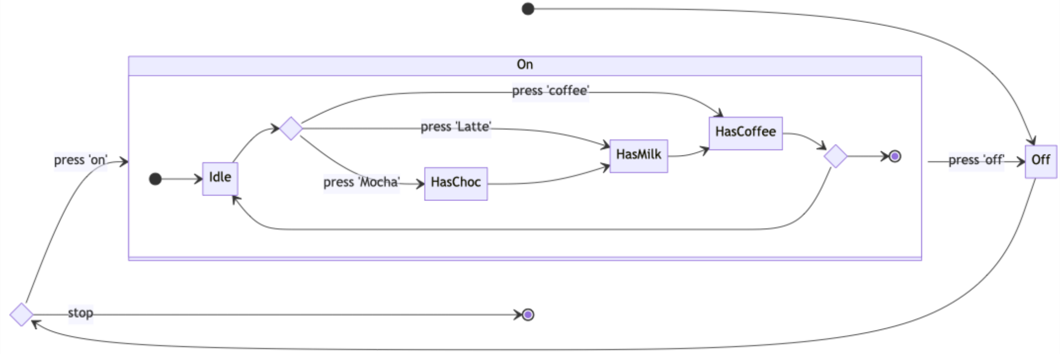
UML behaviour diagrams

UML behaviour diagrams

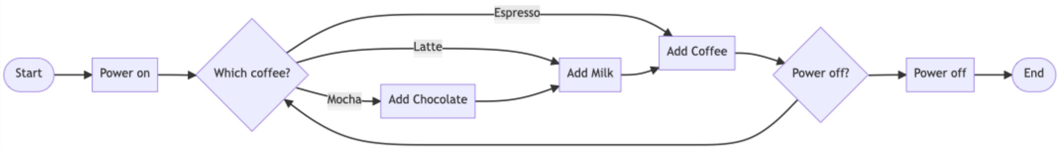
Describe the state of a component, what actions it can do, and how it evolves during its life cycle.

- State Diagram focus on states
- Flowchart focus on actions (also known as *activity diagrams*)

Coffee State Diagram



Coffee Flowchart



Used symbols: *processes*, *decisions*, and *start/end*

Other symbols include: data (or input/output), documents, connectors, comments

Automata – Basic definitions

Sequential and Reactive systems

Sequential systems

Meaning is defined by the results of finite computations

Reactive systems

Meaning is determined by interaction and mobility of non-terminating processes, evolving concurrently

We start here. . .

then we go reactive

Non-Deterministic Finite Automata (NFA)

Definition

A NFA over a set N of names is a tuple $\langle S, I, \downarrow, N, \longrightarrow \rangle$ where

- $S = \{s_0, s_1, s_2, \ldots\}$ is a set of states
- $I \subseteq S$ is the set of initial states
- $\downarrow \subseteq S$ is the set of terminating or final states

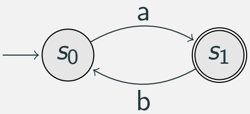
$$\downarrow s \;\equiv\; s \in \downarrow$$

- $\longrightarrow \;\subseteq S \times N \times S$ is the transition relation, often given as an N-indexed family of binary relations

$$s \xrightarrow{a} s' \;\equiv\; \langle s, a, s' \rangle \in \longrightarrow$$

Example

Example of an automaton



s_0 is an initial state

s_1 is a final state

(Formalise this automata)

Ex. 7.1: Formalise these automata as $\langle S, I, \downarrow, N, \longrightarrow \rangle$

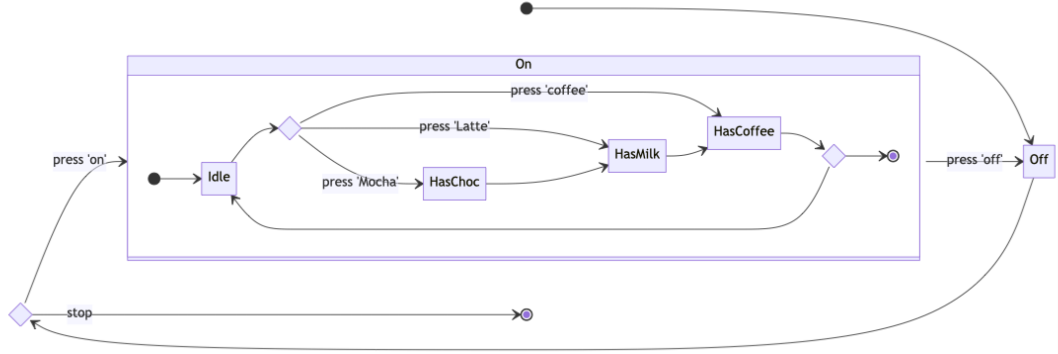
A note on Homework

- 10% of the final mark
- focus on effort – doing badly is better than not doing
- submission: a PDF by email to the teacher who provided the exercises; here pro@isep.ipp.pt.

Deadlines

Exercises presented in a given week must be submitted by the end of the following week, Sunday @ 23h59.
Website/Teams will be kept up-to-date with ongoing open submissions.

Exercise



Ex. 7.2: Draw LTS

(suggestion: by hand on a paper, and take a photo of it.)

Labelled Transition System

More generally, a NFA $\langle S, I, \downarrow, N, \longrightarrow\rangle$ is a labelled transition system (LTS) $\langle S, N, \longrightarrow\rangle$, where each state $s \in S$ determines a system over all states reachable from s and the corresponding restriction of $\longrightarrow$.

LTS classification

- deterministic
- non deterministic
- finite
- finitely branching
- image finite
- ...

Reachability

Definition

The reachability relation, $\longrightarrow^* \subseteq S \times N^* \times S$, is defined inductively

- $s \xrightarrow{\epsilon}{}^* s$ for each $s \in S$, where $\epsilon \in N^*$ denotes the empty word;
- if $s \xrightarrow{a} s''$ and $s'' \xrightarrow{\sigma}{}^* s'$ then $s \xrightarrow{a\sigma}{}^* s'$, for $a \in N, \sigma \in N^*$

Reachable state

$t \in S$ is reachable from $s \in S$ iff there is a word $\sigma \in N^*$ st $s \xrightarrow{\sigma}{}^* t$

Language of an Automaton

Language

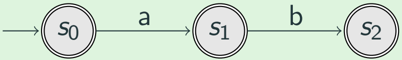
A word σ is in the language L_A of an automata $A = \langle S, I, \downarrow, N, \longrightarrow \rangle$

iff

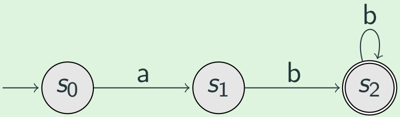
there are states $s \in I$, $s' \in \downarrow$ such that $s \xrightarrow{\sigma}{}^{*} s'$.

Exercises

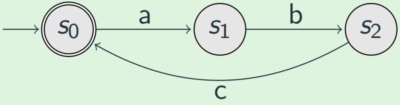
Ex. 7.3: What is the language of this automata?



Ex. 7.4: What is the language of this automata?



Ex. 7.5: What is the language of this automata?



Extra: Regular Expressions

Regular Expressions – syntax

- $w_1 w_2$: word w_1 followed by word w_2
- $w_1 + w_2$: word w_1 or word w_2
- a^*: 0 or more a's
- a^+: 1 or more a's
- ϵ: empty word

Examples

- $ab + c$: (a followed by b) or c
- $(ab)^* b$: b or abb or $ababb$ or ...
- $c((ab)^* b)^+$: cb or $cabb$ or $cababb$ or ...

Extra: Regular Expressions

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Examples

- $ab + c$: (a followed by b) or c
- $(ab)^* b$: b or abb or $ababb$ or ...
- $c((ab)^* b)^+$: cb or $cabb$ or $cababb$ or ...

NFA vs. Reg. Expr.

Word w expressible by a NFA $\Leftrightarrow$ w expressible by a Reg. Expr.

Process algebra

Sequential CCS - Syntax

$$\mathcal{P} \ni P, Q ::= K \mid \alpha.P \mid P + Q \mid \mathbf{0} \mid P[f] \mid P\backslash L \mid P|Q$$

where

- $\alpha \in N \cup \{\tau\}$ is an action
- K s a collection of process names or process constants
- $L \subseteq N$ is a set of labels
- f is a function that renames actions s.t. $f(\tau) = \tau$
- notation:

  $[f] = [a_1 \mapsto b_1, \ldots, a_n \mapsto b_n]$

Syntax

$$\mathcal{P} \ni P, Q ::= K \mid \alpha.P \mid P + Q \mid \mathbf{0} \mid P[f] \mid P \backslash L \mid P|Q$$

Ex. 7.6: Which are NOT syntactically correct? Why?

$$a.b.A + B \quad (1)$$
$$(a.\mathbf{0} + b.A) \backslash \{a, b, c\} \quad (2)$$
$$(a.\mathbf{0} + b.A) \backslash \{a, \tau\} \quad (3)$$
$$a.B + [b \mapsto a] \quad (4)$$
$$\tau.\tau.B + \mathbf{0} \quad (5)$$
$$a.(a + b).A \quad (6)$$
$$(a.B + b.B)[a \mapsto a, \tau \mapsto b] \quad (7)$$
$$(a.B + \tau.B)[b \mapsto a, a \mapsto a] \quad (8)$$
$$(a.b.A + b.\mathbf{0}).B \quad (9)$$
$$(a.b.A + b.\mathbf{0}) + B \quad (10)$$

CCS semantics - building a NFA

$$\text{(act)}\quad \frac{}{\alpha.P \xrightarrow{\alpha} P}$$

$$\text{(sum-1)}\quad \frac{P_1 \xrightarrow{\alpha} P_1'}{P_1 + P_2 \xrightarrow{\alpha} P_1'}$$

$$\text{(sum-2)}\quad \frac{P_2 \xrightarrow{\alpha} P_2'}{P_1 + P_2 \xrightarrow{\alpha} P_2'}$$

$$\text{(res)}\quad \frac{P \xrightarrow{\alpha} P'}{P \backslash L \xrightarrow{\alpha} P' \backslash L}\ \alpha \notin L$$

$$\text{(rel)}\quad \frac{P \xrightarrow{\alpha} P'}{P[f] \xrightarrow{f(\alpha)} P'[f]}$$

- Initial states: the process being translated
- Final states: all states are final
- Language: possible sequence of actions of a process

CCS semantics - building a NFA

$$\text{(act)}\quad \frac{}{\alpha.P \xrightarrow{\alpha} P}$$

$$\text{(sum-1)}\quad \frac{P_1 \xrightarrow{\alpha} P_1'}{P_1 + P_2 \xrightarrow{\alpha} P_1'}$$

$$\text{(sum-2)}\quad \frac{P_2 \xrightarrow{\alpha} P_2'}{P_1 + P_2 \xrightarrow{\alpha} P_2'}$$

$$\text{(res)}\quad \frac{P \xrightarrow{\alpha} P'}{P\backslash L \xrightarrow{\alpha} P'\backslash L}\ \alpha \notin L$$

$$\text{(rel)}\quad \frac{P \xrightarrow{\alpha} P'}{P[f] \xrightarrow{f(\alpha)} P'[f]}$$

Ex. 7.7: Build a derivation tree to prove the transitions below

1. $(a.A + b.B) \xrightarrow{b} B$
2. $(a.b.A + (b.a.B + c.a.C)) \xrightarrow{b} a.B$
3. $((a.B + b.A)[a \mapsto c])\backslash\{a, b\} \xrightarrow{c} (B[a \mapsto c])\backslash\{a, b\}$

Exercise

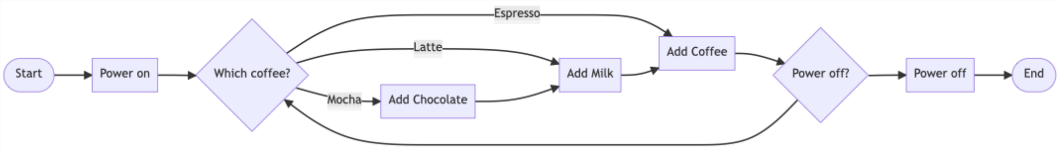
Ex. 7.8: Draw the automata

$$CM = \text{coin.coffee.}CM$$
$$CS = \text{pub.(coin.coffee.CS} + \text{coin.tea.CS)}$$

Ex. 7.9: What is the language of the process A?

$$A = \text{goLeft.A} + \text{goRight.B}$$
$$B = \text{rest.}\mathbf{0}$$

Exercise



Ex. 7.10: Write the process of the flowchart above

$$
\begin{aligned}
P &= \mathsf{powerOn}.Q \\
Q &= \mathsf{selMocha}.\mathsf{addChocolate}.Mk + \mathsf{selLatte}.Mk + \ldots \\
Mk &= \mathsf{addMilk}\ldots
\end{aligned}
$$

mCRL2 Tools

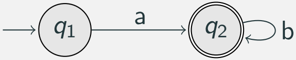
Slides 10:

`https://cister-labs.github.io/ramde2122/slides/10-mcrl2.pdf`

Concurrent Process algebra

Overview

Recall

1. Non-deterministic Finite Automata: 

2. (Sequential) Process algebra: $P = a.Q \quad Q = b.Q$
3. Meaning of (2) using (1)

Still missing

- **Interaction** between processes
- *Interaction diagrams* vs. *interacting processes*
- Enrich (2) and (3)

CCS - Updated Syntax

$$\mathcal{P} \ni P, Q ::= K \mid \alpha.P \mid P + Q \mid \mathbf{0} \mid P[f] \mid P\backslash L \mid P|Q$$

where

- $\alpha \in N \cup \overline{N} \cup \{\tau\}$ is an action
- K s a collection of process names or process constants
- $L \subseteq N$ is a set of labels
- f is a function that renames actions s.t. $f(\tau) = \tau$ and $f(\bar{a}) = \overline{f(a)}$
- notation:

  $[f] = [a_1 \mapsto b_1, \ldots, a_n \mapsto b_n]$ where $a_i, b_i \in N \cup \{\tau\}$

Syntax

$$\mathcal{P} \ni P, Q ::= K \mid \alpha.P \mid P + Q \mid \mathbf{0} \mid P[f] \mid P\backslash L \mid P|Q$$

Ex. 7.11: Which are syntactically correct?

$$a.\overline{b}.A + B \quad (11)$$
$$(a.\mathbf{0} + \overline{a}.A)\backslash\{\overline{a}, b\} \quad (12)$$
$$(a.\mathbf{0} + \overline{a}.A)\backslash\{a, \tau\} \quad (13)$$
$$(a.\mathbf{0} + \overline{\tau}.A)\backslash\{a\} \quad (14)$$
$$\tau.\tau.B + \overline{a}.\mathbf{0} \quad (15)$$
$$(\mathbf{0}|\mathbf{0}) + \mathbf{0} \quad (16)$$
$$(a.B + b.B)[a \mapsto a, \tau \mapsto b] \quad (17)$$
$$(a.B + \tau.B)[b \mapsto a, b \mapsto a] \quad (18)$$
$$(a.B + b.B)[a \mapsto b, b \mapsto \overline{a}] \quad (19)$$
$$(a.b.A + \overline{a}.\mathbf{0})|B \quad (20)$$
$$(a.b.A + \overline{a}.\mathbf{0}).B \quad (21)$$
$$(a.b.A + \overline{a}.\mathbf{0}) + B \quad (22)$$

CCS semantics - building an NFA

(act)
$$\frac{}{\alpha.P \xrightarrow{\alpha} P}$$

(sum-1)
$$\frac{P_1 \xrightarrow{\alpha} P_1'}{P_1 + P_2 \xrightarrow{\alpha} P_1'}$$

(sum-2)
$$\frac{P_2 \xrightarrow{\alpha} P_2'}{P_1 + P_2 \xrightarrow{\alpha} P_2'}$$

(res)
$$\frac{P \xrightarrow{\alpha} P'}{P\backslash L \xrightarrow{\alpha} P'\backslash L} \quad \alpha, \overline{\alpha} \notin L$$

(rel)
$$\frac{P \xrightarrow{\alpha} P'}{P[f] \xrightarrow{f(\alpha)} P'[f]}$$

(com1)
$$\frac{P \xrightarrow{\alpha} P'}{P|Q \xrightarrow{\alpha} P'|Q}$$

(com2)
$$\frac{Q \xrightarrow{\alpha} Q'}{P|Q \xrightarrow{\alpha} P|Q'}$$

(com3)
$$\frac{P \xrightarrow{a} P' \quad Q \xrightarrow{\overline{a}} Q'}{P|Q \xrightarrow{\tau} P'|Q'}$$

CCS semantics - building an NFA

$$\text{(act)}\quad \frac{}{\alpha.P \xrightarrow{\alpha} P} \qquad \text{(sum-1)}\quad \frac{P_1 \xrightarrow{\alpha} P_1'}{P_1 + P_2 \xrightarrow{\alpha} P_1'} \qquad \text{(sum-2)}\quad \frac{P_2 \xrightarrow{\alpha} P_2'}{P_1 + P_2 \xrightarrow{\alpha} P_2'}$$

$$\text{(res)}\quad \frac{P \xrightarrow{\alpha} P'}{P\backslash L \xrightarrow{\alpha} P'\backslash L}\ \alpha, \overline{\alpha} \notin L \qquad \text{(rel)}\quad \frac{P \xrightarrow{\alpha} P'}{P[f] \xrightarrow{f(\alpha)} P'[f]}$$

$$\text{(com1)}\quad \frac{P \xrightarrow{\alpha} P'}{P|Q \xrightarrow{\alpha} P'|Q} \qquad \text{(com2)}\quad \frac{Q \xrightarrow{\alpha} Q'}{P|Q \xrightarrow{\alpha} P|Q'} \qquad \text{(com3)}\quad \frac{P \xrightarrow{a} P' \quad Q \xrightarrow{\overline{a}} Q'}{P|Q \xrightarrow{\tau} P'|Q'}$$

Ex. 7.12: Draw the NFAs

$$CM = \text{coin}.\overline{\text{coffee}}.CM$$

$$CS = \text{pub}.\overline{\text{coin}}.\text{coffee}.CS$$

$$SmUni = (CM|CS)\backslash\{\text{coin}, \text{coffee}\}$$

Exercises

Ex. 7.13: Let $A = b.a.B$. Show that:

1. $(A \mid \overline{b}.\mathbf{0})\backslash\{b\} \xrightarrow{\tau} (a.B \mid \mathbf{0})\backslash\{b\}$
2. $(A \mid b.a.B) + ((b.A)[b \mapsto a]) \xrightarrow{a} A[b \mapsto a]$

Ex. 7.14: Draw the NFAs A and D

$$A = x.B + x.x.C \qquad D = x.x.x.D + x.E$$

$$B = x.x.A + y.C \qquad E = x.F + y.F$$

$$C = x.A \qquad F = x.A$$

mCRL2 Tools

Slides 10:

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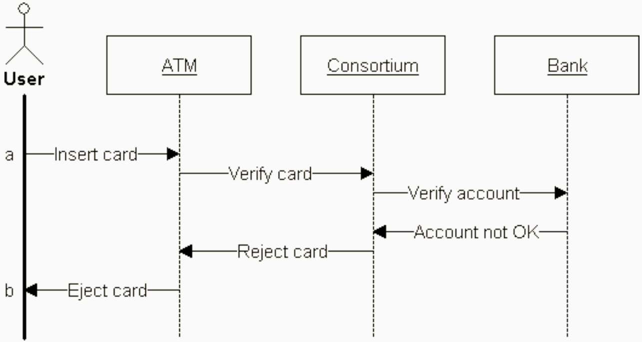
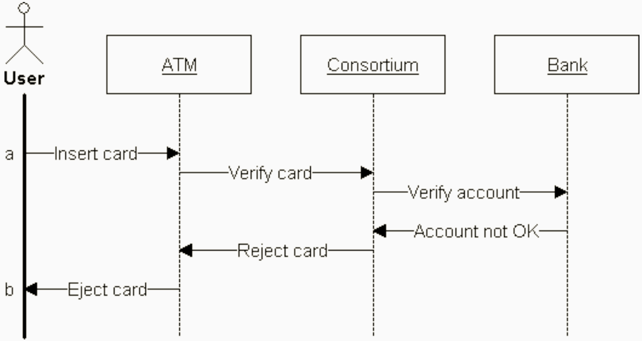
Sequence Diagrams vs. Interactive Processes

Sequence Diagrams as Interactive Processes



- Objects as Processes (e.g., processes U, A, C, B)
- Send actions (e.g., $\overline{insertCard}$)
- Reveive actions (e.g., $insertCard$)

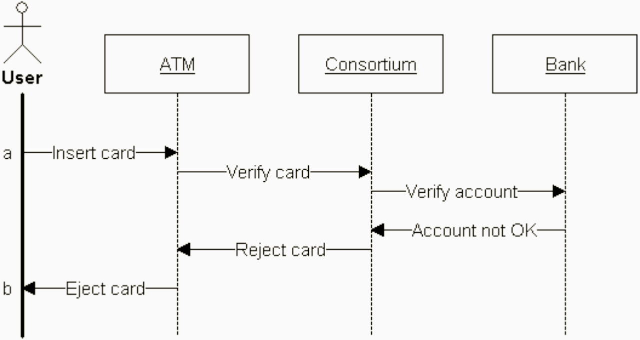
Language of Sequence Diagrams, Informally



This example has only 1 word and its prefixes

$$Sys_{global} = \mathit{insertCard} \; . \; \mathit{verifyCard} \; . \; \mathit{verifyAccount} \; . \; \mathit{accountNotOK} \; . \; \mathit{rejectedCard} \; . \; \mathit{ejectCard} \; . \; \mathbf{0}$$

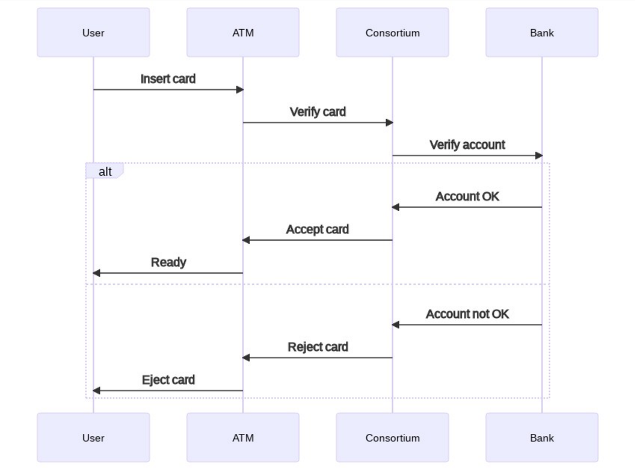
Sequence Diagrams as Interactive Processes



Ex. 7.15: Write an interactive processes that act as the seq. diagram.

$$Sys_{local} = (U|A|C|E)\backslash \ldots$$
$$U = \overline{insertCard}.ejectCart.\mathbf{0}$$
$$A = \ldots$$
$$C = \ldots$$
$$E = \ldots$$

Sequence Diagrams as Interactive Processes



Ex. 7.16: Write a single process Sys_{global} and a set of interactive processes Sys_{local} that act as the seq. diagram.

$$Sys_{global} = insertCard\ .\ \ldots$$

$$Sys_{local} = (U|A|C|E)\backslash \ldots$$

$$U = \ldots$$

$$A = \ldots$$

$$C = \ldots$$

$$E = \ldots$$