

RAMDE

Requirements and Model Drive Engineering

Introducing Formal Languages & Methods
in Requirements Engineering



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What we will be looking at now

- Why **do we need** Formal Methods in Requirements Engineering
- What do Formal Methods **have to offer**

“It is clear to all the best minds in the field that a more mathematical approach is needed for software to progress much.”

Bertand Meyer, author of the concept of "Design by Contract"

What are Formal Methods?

A very, very lightweight introduction...



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What are Formal Methods?

➤ General view:

- application of **mathematics** to **software engineering**
- concerned with **specifying**, **modelling**, and **analysing** the system using and underlying, **precise**, **mathematically based** notation

➤ More narrow view:

- use **formal languages**
- engage on **formal reasoning** about **formulae** of the language

Formal Methods and Requirements

- How do Formal Methods relate to requirements engineering?
- Essentially, through **formal notation**:
 - Formal **set of rules** that define the **syntax** and **semantics** of the language
 - The rules that determine if the formulae are **well-formed** and to **prove properties** about those formulae

Highlight Example

➤ First Order Logic (also known as Predicate Logic)

- primitives for **expressions**: variables, constants, functions, predicates
- logical **connectives**: and ($\wedge$), or ($\vee$), implies ($\rightarrow$), equality ($=$), negation ($\neg$)
- **quantifiers**: for all ($\forall$), exists ($\exists$)
- **deduction** rules: we will look into this in the next lessons

➤ Examples of expression/formulae:

- $(x > y) \wedge (y > z) \rightarrow x > z$
- $\forall x, \forall y, \forall z, (x > y) \wedge (y > z) \rightarrow x > z$
- $\forall x, \forall y, \forall z, [\text{Node}(x, y, z) \rightarrow \text{Leaf}(x) \rightarrow \text{Leaf}(z)] \rightarrow x < y \wedge z \geq y$
- $\forall h, \forall l, \text{List}(l) \rightarrow \text{List}(h :: l)$
- $\forall p, \forall s, \text{Parent}(p, s) \rightarrow \neg \text{Parent}(s, p)$

Applying Formal Methods

Understanding how domain,
requirements, and specification relate



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What does correctness mean?

➤ Let us consider the following representation



➤ Some important information:

- *Domain properties* are **characteristics** of the **application domain**
- *Requirements* are **conditions/features** that to be **made true** by the system
- *Specification* is a **description of the behaviours** the **system must have** in order to **satisfy the requirements**

What does correctness mean?

› Correctness criteria:

- › the program running on a given computer **satisfies** the **specification**
- › the **specification**, in the context of the domain properties, **satisfies** the requirements

› Completeness criteria:

- › We **discovered** all the **important requirements**
- › We **discovered** all the **relevant domain properties**

Understanding the Differences

> Requirement **R**:

The database shall only be accessible by authorized personnel

> Domain Properties **D**:

Authorized personnel have passwords

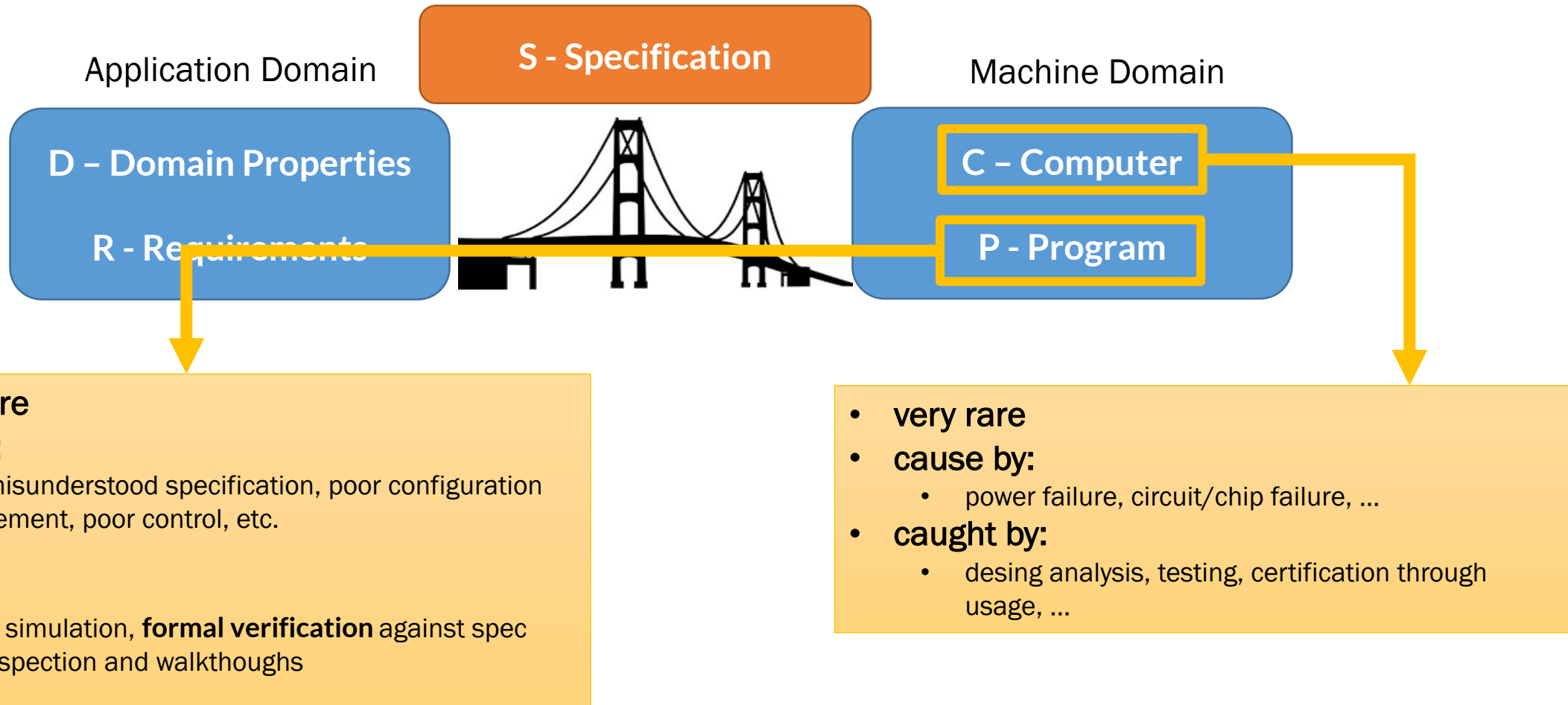
Passwords are never shared with non-authorized personnel

> Specification **S**:

*Access to the database shall only be granted after the user types
and authorized password*

D + S implies R

When things go wrong



When things go wrong



- common
- caused by:
 - misunderstood requirements
 - wrong choice of specification language
 - ambiguous, inconsistent/incomplete specification
- caught by:
 - inspection, **formal verification**
 - end-to-end testing

When things go wrong

Application Domain

S - Specification

Machine Domain

D - Domain Properties

R - Requirements

C - Computer

P - Program

- **very common**
- **caused by:**
 - lack of domain expertise
 - unquestioned assumptions
 - insufficient domain analysis
- **caught by:**
 - failure analysis
 - talking to the right experts
 - ...

- **common**
- **caused by:**
 - insufficient/bad communication with stakeholders
 - lack of analysis
 - failure to handle change
- **caught by:**
 - inspections, customer reviews
 - modelling, **formal verification**, prototyping

Where do Formal Methods Apply

➤ Formalize the specification:

- as a precise baseline to **verify** the **program** against
- as a **model of program behavior** to compare against **requirements**

➤ Formalize the domain knowledge:

- **reason** about whether it is **complete**
- **reason** about how it **affects** the proposed system

➤ Formalize the requirements:

- **animate** them
- test **logical coherence**
- check for **completeness against** an underlying mathematical **model**

Why formalise in Requirements Eng?

- To remove ambiguity and improve precision
- Provides a basis for verification:
 - That a **program** meets its **specification**
 - That a that **specification captures** the **requirements** adequately
- Allows us to reason about the requirements
 - **Properties** of formal requirements models can be **checked automatically**
 - Can test for **consistency**, explore the **consequences**, etc.
- Allows us to animate/execute the requirements
- Will have to formalize eventually anyway
 - **bridging** from the informal **world** to a formal **machine** domain

Why is not usually done?

- Tend to exhibit lower level than other techniques
 - inclusion of a **lot of detail** and the **system boundaries** already well **established**
- Tend to concentrate on consistent, correct models
- Confusion about which tools are appropriate
 - advocates generally **too tied** to a **certain tool**, and **scalability** is still **not possible** with many tools/prototypes
- Require more effort
 - mostly due to the necessary **mathematical training** and **experience**
- Are not appropriate in many projects...

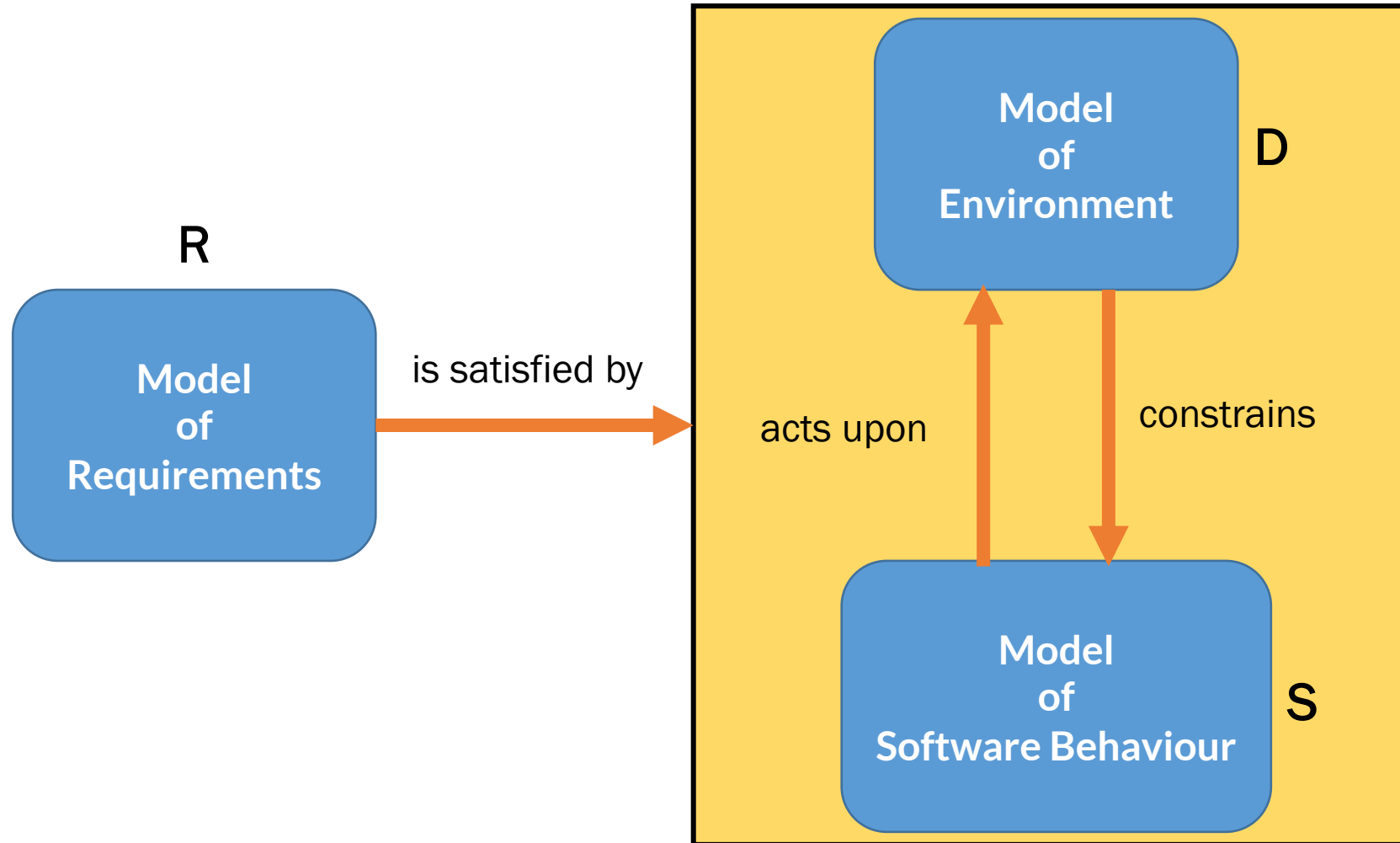
Modelling

And how it relates to Formal Methods



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So, too many models involved?



Modelling can help!

- Can guide elicitation:

- understanding what **questions** to make, **identify hidden** requirements

- Can provide a measure of progress:

- understand if **completeness** of the model implies **finished elicitation**

- Uncover problems:

- **inconsistency** of the model **revealing** other interesting **aspects** of the system (e.g., *infeasible requirements, terminology confusion, disagreement between stakeholders*)

- Check our own understanding:

- model **has the right** properties, or **reason** about its **consequences**

Types of Model

> Informal, written in natural language

- > extremely **expressive** and **flexible**, **poor** at capturing the **semantics** of the model, good for **elicitation** and to **annotate** models for communication

> Semi-formal

- > captures **structure** and **some semantics**
- > allow for some **reasoning**, **consistency** checking, etc.

> Formal

- > very **precise semantics**, possible to conduct **reasoning**

Good Properties of Modelling Languages

› Independent of Implementation

- › **does not** consider **internal representation**, abstracts data types

› Abstraction capabilities

- › extracts the **essential aspects** of the system, allowing to have a “big pic”

› Formality

- › precise **syntax** and rich **semantics**

› Constructability/Composability

- › can **compose/decompose** pieces of the model to handle **complexity** and size

Good Properties of Modelling Languages

➤ Ease of analysis

- ability to **analyze** for **ambiguity, incompleteness, inconsistency**

➤ Traceability

- ability to **cross-reference elements**, ability to **link** to **design**, etc.

➤ Executability

- can **animate** the model, to **compare** it to **reality**

➤ Minimality

- no **redundancy** of **concepts**, **no extraneous choices** of how to represent **concepts**

Model Validation

- **Consistency Analysis and Typechecking**
 - **well-formedness first** that most of other methods and represent real-world integrity
- **Validation**
 - **animation** of small models
 - **formal queries** about **implications** of know properties, particular requirements, or changes
 - **system properties**
- **Verifying design refinement**
 - does the **design meets the requirements**

Some differences between Formal Methods

Formal methods **differ considerably** on their nature, notably in what concerns the following aspects.

- Ontology
 - **fixed** vs **extensible**
- Mathematical foundations
 - **logic** vs **algebraic languages** vs **set theory** vs **state machines**
- Treatment of time
 - **discrete sequences of events** vs **continuous time** vs **quantified intervals**

Some Traditions/Schools

➤ Formal Specification Languages

- inception on **program verification**, suitable for **specifying** the **behaviour** of **program units**
- key technologies being **type checking**, **theorem proving**

➤ Reactive System Modelling

- capture **dynamic models** of **system behaviour**, with focus on **reactive systems** (e.g., safety and liveness in real-time embedded systems)
- key technologies: **model checking**, **consistency checking**

➤ Formal Conceptual Modeling

- focus on **modelling domain entities**, **activities**, **agents**, **assertions**
- key technologies: **formal ontologies**, **first order logic**

Formal Specification Languages

› Three families:

- › **Operational** - specification is executable abstraction of the implementation
- › **State-based** - views a program as a (large) data structures whose state can be altered by procedure calls
- › **Algebraic** - views a program as a set of abstract data structures with a set of operations

› Developed for specifying programs

- › **Programs** are **formal**, man-made **objects**
- › these languages are typically **not very appropriate** for Requirements Engineering (**requirements specification $\neq$ program specification**)

Reactive Systems Modelling

- Modeling how a system should behave
 - model the **environment** as a **state machine**
 - model the **system** as a **state machine**
 - model **safety**, **liveness** properties of the machine as **temporal logic assertions**
 - check whether the **properties hold** of the system **interacting** with its **environment**
- We will look into this during the course, mostly in the 2nd semester
 - class on formal verification of software
 - in the current class you will get the basis using algebraic specifications

From Notations to Methods

> Notation:

- > a representation **language** for **expressing things**

> Technique:

- > prescribes **how to perform** a particular activity and **its product** in a **particular notation**

> Method:

- > definition on **how to perform** a collection of **activities**, focusing on **integration of techniques** and **guidance** about their use

Conclusions

- **Formal languages and methods** are here to **help** in the very **complex and challenging** process of Requirement Engineering
- **Requirements** valid as a **consequence** of **domain properties** and **software properties & behaviour**
- **Several flavours** of formal methods available, which we must **choose wisely**
- **First Order Logic as highlight method**
 - *we will start looking at it in the next lesson, picking up on the basic mathematical notations, and proceeding to syntax, semantics, and deductive rules that support reasoning*

The End



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