

Natural Deduction Rules for Propositional Logic

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Conjunction

The rules:

$$\frac{\varphi \quad \psi}{\varphi \wedge \psi} \wedge \mathbf{I} \qquad \frac{\varphi \wedge \psi}{\varphi} \wedge \mathbf{E}_l \qquad \frac{\varphi \wedge \psi}{\psi} \wedge \mathbf{E}_r$$

Representation in Fitch's style:

$$\begin{array}{c} \vdots \\ n \quad | \quad \varphi \\ \vdots \\ m \quad | \quad \psi \\ \vdots \\ \varphi \wedge \psi \end{array} \wedge \mathbf{I}(n, m) \qquad \begin{array}{c} \vdots \\ n \quad | \quad \varphi \wedge \psi \\ \vdots \\ m \quad | \quad \varphi \end{array} \wedge \mathbf{E}_l(n) \qquad \begin{array}{c} \vdots \\ n \quad | \quad \varphi \wedge \psi \\ \vdots \\ m \quad | \quad \psi \end{array} \wedge \mathbf{E}_r(n)$$

Disjunction

The rules:

$$\frac{\varphi}{\varphi \vee \psi} \vee \mathbf{I}_l \qquad \frac{\psi}{\varphi \vee \psi} \vee \mathbf{I}_r \qquad \frac{\begin{array}{c} [\varphi] \\ \vdots \\ \theta \end{array} \quad \begin{array}{c} [\psi] \\ \vdots \\ \theta \end{array}}{\varphi \vee \psi \quad \theta} \vee \mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c}
 n \quad \left| \begin{array}{l} \vdots \\ \varphi \\ \vdots \\ \varphi \vee \psi \end{array} \right. \quad \forall \mathbf{I}_l(n)
 \end{array}
 \quad
 \begin{array}{c}
 n \quad \left| \begin{array}{l} \vdots \\ \psi \\ \vdots \\ \varphi \vee \psi \end{array} \right. \quad \forall \mathbf{I}_r(n)
 \end{array}
 \quad
 \begin{array}{c}
 \vdots \\
 n \quad \varphi \vee \psi \\
 \vdots \\
 m_1 \quad \left| \begin{array}{l} \varphi \\ \vdots \\ \theta \end{array} \right. \\
 k_1 \\
 \vdots \\
 m_2 \quad \left| \begin{array}{l} \psi \\ \vdots \\ \theta \end{array} \right. \\
 k_2 \\
 \vdots \\
 \theta
 \end{array}
 \quad
 \forall \mathbf{E}(n, m_1 - k_1, m_2 - k_2)$$

Negation

The rules:

$$\begin{array}{c}
 [\varphi] \\
 \vdots \\
 \perp \\
 \hline
 \neg \varphi \quad \neg \mathbf{I}
 \end{array}
 \quad
 \frac{\neg \neg \varphi}{\varphi} \neg \mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c}
 n \\
 \vdots \\
 \hline
 \varphi \\
 \vdots \\
 m \\
 \hline
 \perp \\
 \hline
 \neg\varphi
 \end{array}
 \quad \neg\mathbf{I}(n-m)$$

$$\begin{array}{c}
 n \\
 \vdots \\
 \neg\neg\varphi \\
 \vdots \\
 \varphi
 \end{array}
 \quad \neg\mathbf{E}(n)$$

False

The rules:

$$\begin{array}{c}
 \varphi \\
 \vdots \\
 \neg\varphi \\
 \hline
 \perp
 \end{array}
 \quad \perp\mathbf{I}$$

$$\frac{\perp}{\varphi} \quad \perp\mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c}
 n \\
 \vdots \\
 \varphi \\
 \vdots \\
 m \\
 \vdots \\
 \neg\varphi \\
 \vdots \\
 \perp
 \end{array}
 \quad \perp\mathbf{I}(n, m)$$

$$\begin{array}{c}
 n \\
 \vdots \\
 \perp \\
 \vdots \\
 \varphi
 \end{array}
 \quad \perp\mathbf{E}(n)$$

Implication

The rules:

$$\begin{array}{c}
 [\varphi] \\
 \vdots \\
 \psi \\
 \hline
 \varphi \rightarrow \psi
 \end{array}
 \quad \rightarrow\mathbf{I}$$

$$\frac{\varphi \rightarrow \psi \quad \varphi}{\psi} \quad \rightarrow\mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{l}
 n \\
 \vdots \\
 \hline
 \varphi \\
 \hline
 \vdots \\
 m \\
 \psi \\
 \varphi \rightarrow \psi \quad \rightarrow\mathbf{I}(n-m)
 \end{array}$$

$$\begin{array}{l}
 \vdots \\
 n \\
 \varphi \rightarrow \psi \\
 \vdots \\
 m \\
 \varphi \\
 \vdots \\
 \psi \quad \rightarrow\mathbf{E}(n, m)
 \end{array}$$